Reversible Phase Transitions in a Structured Overlay Network with Churn

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<u>Introduction</u>

- Applications are exposed to increasingly stressful environments
 - Out of data centers, to georeplication and edge computing
 - Node and communication failures are increasing as nodes increase in number
- We would like applications to survive such stressful environments and to have predictable behavior
 - We introduce the concept of Reversibility to define what survival means in arbitrarily stressful environments
 - We introduce the concept of Phase to allow applications to observe the hostility of their environment and behave accordingly
- We evaluate these concepts on a large realistic system
 - A structured overlay network with simulated environment and high churn
 - We investigate how to make it Reversible and how to build applications on top

Reversibility

Why we need Reversibility

- Suppose a distributed system running on n nodes providing a specific set of services
- From time t to t+T, the system experiences external stress
 - Ex. k nodes crash and j nodes join the system
 - Ex. a system partition due to a connectivity problem of the underlying physical network
- Can we ensure that the system will eventually regain its full functionality after time t+T?
- Does the system have a well-defined behavior during the interval [t,t+T]?

Reversibility (informal)

 With Reversibility we can give affirmative answers to both questions!

 Informally, Reversibility means that the system's functionality depends only on the current stress experienced by the system and not on the history of the stress

Reversibility (formal)

- Given a function S(t) that returns the system stress in some arbitrary but well-defined units
 - Ex. S(t) can explain how the system is partitioned as a function of time,
 or give churn as function of time

- A system is **Reversible** if there exists a function $F_{op}(id,S(t))$ of node identifier id and stress S(t) such that the set of system operations available at node id is $F_{op}(id,S(t))$
 - An operation is available for a given stress if the operation will eventually succeed (it will fail only a finite number of times if tried repeatedly and then succeed)
 - Note that when S(t)=0 the system provides full functionality

Comparison with related concepts

Reversibility versus Fault Tolerance

- A fault-tolerant system is resilient for a given fault model, but its behavior outside that model is undefined
- Reversibility is a stronger property because it guarantees that the system will recover functionality if the stress is removed

Reversibility versus Self Stabilization

- A self-stabilizing system survives any temporary perturbation of its internal state; it returns to a valid state when there are no perturbations
- Reversibility is more useful in practice: it gives information about functionality even during nonzero stress

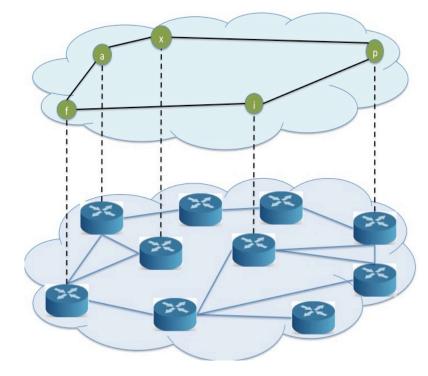
Evaluation

Evaluation

- Investigate Reversibility in the context of a realistic system
 - Representative system: a structured overlay network
- Simulated environment running on Mozart-Oz 2.0 platform
 - Simulated message delays follow Internet distribution
 - Network size of 1024 peers
- Experiments
 - First story: achieving Reversibility during high churn
 - Second story: deducing system functionality by observing structure
 - Third story: designing Reversible applications

Structured Overlay Network (SON)

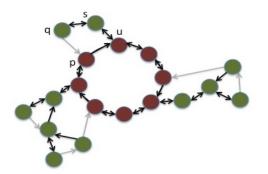
- P2P Systems: Dual client/server role of each node of the system.
- Due to local cooperation of peers an overall network routing view emerges, known as an *overlay* network, on top of the underlay network.
- Structured Overlay Network: A structure is induced through the pointers maintained by each peer of the system.



Overlay Network: A P2P System with nodes a, f, i, p and x forms the overlay network on top of the underlay network

<u>Beernet</u>

- Beernet³ is a representative example of the design class as per the reference architecture proposed by Aberer et.al.
- Why Beernet?
 - Similar to Chord, but with correct lock-free join operation.
 - Join/leave in Chord requires coordination of three peers that is not guaranteed due to non-transitive connectivity on Internet.
 - Non-Transitive Connectivity: A can talk to B and B can talk to $C \neq A$ can talk to C.
 - Beernet does not assume transitive connectivity. More resilient on Internet. Three step join/leave operation, each step requires coordination among only two peers (guaranteed with a point-to-point communication).
 - Consequence: Natural Branching structure. A stable core ring and transient branches.



Branches on a relaxed ring. Peers p and s consider u as successor, but u only considers s as predecessor. Peer q has not established a connection with its predecessor p yet.

Maintenance Strategies

- A Maintenance Strategy maintains correct structure of a SON
 - We investigate the Maintenance Strategies needed for Reversibility
- Several strategies are proposed in the literature:
 - Correction-on-Change/Use (used by DKS, Beernet);
 - Periodic Stabilization (used by Chord);
 - Gossip-based strategies, e.g., T-MAN (building overlay topology).
- These strategies are complementary
 - Correction-on-change is much more efficient than gossip, whereas gossip is much more resilient

Maintenance Strategies (cont..)

• We cover a complete space of possible maintenance strategies:

Efficiency	Maintenance Strategy	Local/ Global	Reactive/ Proactive	Fast/ Slow	Safety	Bandwidth Consumption
+	Correction-on-Change (for self-healing) and Correction-on-Use (provides self-optimization and self-configuration).	Local	Reactive	Fast	Yes	Small
+	Periodic Stabilization: correction using periodic probing.	Local	Proactive	Slow	Lookup inconsistencies and uncorrected false suspicions can be introduced	High
+	Overlay Merger with Passive List: Trigger Merger using falsely suspected nodes ²	Global	Reactive	Adaptable	Yes	Adaptable
	Gossip-based Maintenance, e.g., Overlay Merger ² with Knowledge Base: Proactive approach to trigger merger using the gathered knowledge at each node.	Global	Proactive	Adaptable	Yes	Adaptable
Resiliency						

Stories and Their Contributions

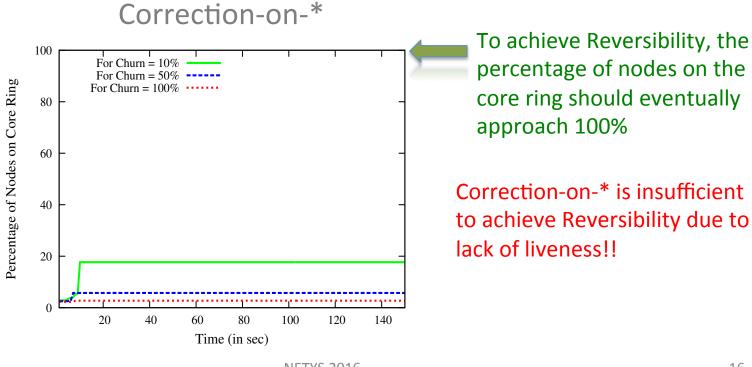
- First Story: "Can the system be made reversible against churn using the Maintenance Strategies?"
 - We show experimentally the need of both efficient and resilient maintenance
- Second Story: "Can we deduce the system's functionality by examining its structure at high churn? YES! Phase concept."
 - Insight on how to observe global structure;
 - Insight on how phase of each node is related to functionality of the system;
 - Experimental demonstration that reversible phase transitions happen in a reversible system as the stress varies
- Third Story: "Can we help applications to be reversible and predictable"?
 YES! Expose Phase of each node through an API."
 - Introduction of Phase API;
 - Insight on how the application can use phase concept to manage its behavior

First Story Churn & Reversibility

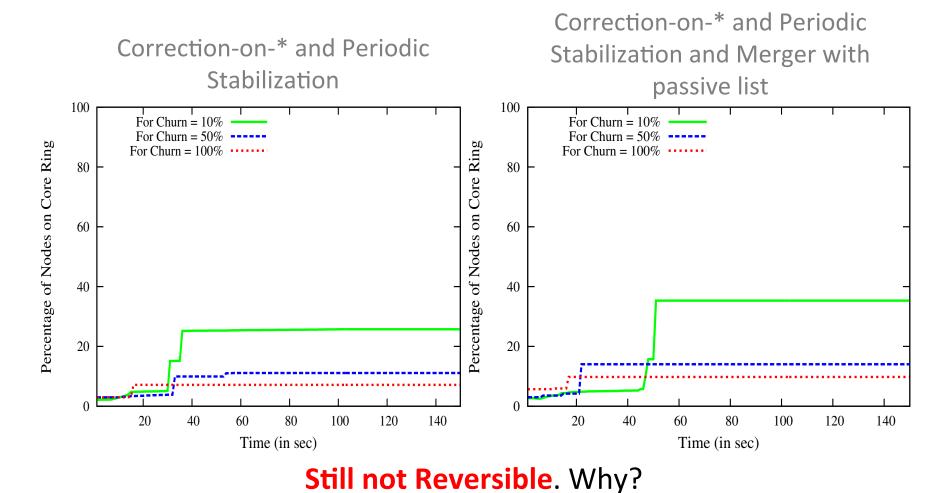
Are the Maintenance Strategies Reversible? (1)

Churn: % of node turnover per second.

Metric: % of nodes on core ring as a function of time

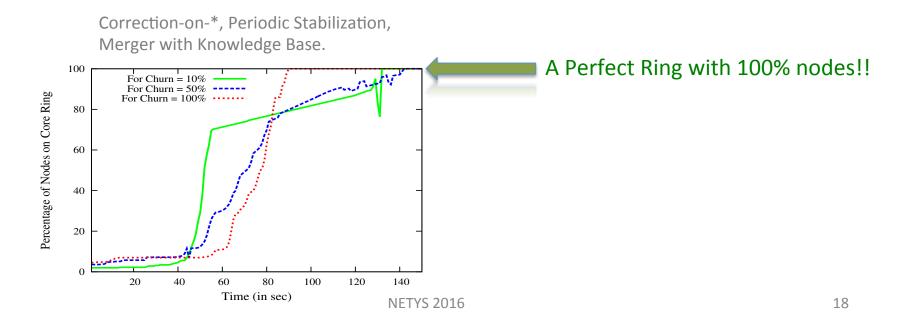


Are the Maintenance Strategies Reversible? (2)



Are the Maintenance Strategies Reversible? (3)

- High churn makes overlay unstable, which does not allow new peers to complete a join
 - The churn rapidly invalidates the join reference of the new peer
- In order to make these isolated peers part of overlay, we need to re-trigger join by providing a new valid join reference.
 - Knowledge Base is required to get knowledge about an alive peer of overlay
- Proactive triggering of merger using Knowledge Base to avoid partition of the system after isolated nodes complete their join procedures.



Summary of First Story

 Repeated join using Knowledge Base is required to achieve Reversibility against extremely high Churn.

 Proactive merger using Knowledge Base is required to avoid partitioning of the system.

Second Story Phase and Phase Transitions

Phase, Phase Transition & Critical Point

- System = An aggregate entity composed of a large number of interacting parts
 - Each part is a node of the SON
- A *Phase* is a subset of a system for which the qualitative properties (e.g., functional guarantees) are essentially the same
 - Different parts can be in different phases, depending on the local environment observed by the part
- Why is this interesting?
 - System functionality depends on these qualitative properties
 - Use phase for observing system functionality, but it should work without extra computation and even when communication is broken
 - Useful to applications running on top of SON in stressful environments

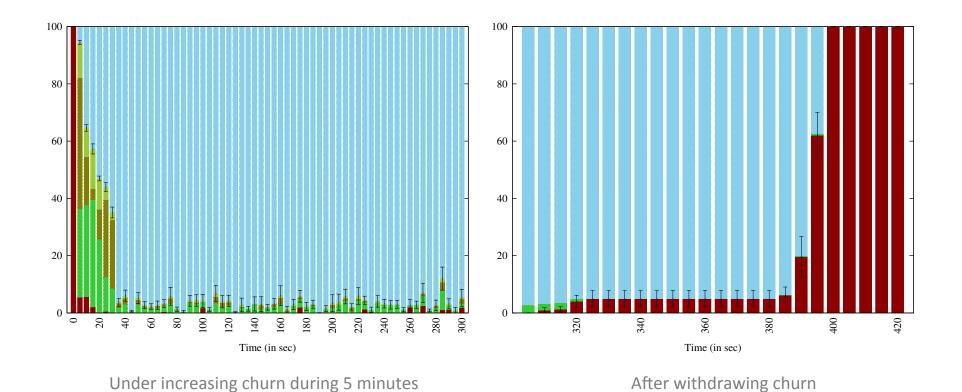
Phase, Phase Transition & Critical Point (Cont..)

- A Phase Transition occurs when a significant fraction of a system's parts changes phase
 - This can happen if the local environment changes at many parts
- A Critical Point occurs when more than one phase exists simultaneously in significant fractions of a system
- Reversibility and Phase:
 - Stress is a global condition that cannot be easily measured by individual nodes
 - Phase P_i at each node i is a well-defined local property
 - Phase configuration of system, $P_c = (P_1, P_2, P_3, ..., P_n)$.
 - The set of available operations of the system, namely F_{det} (id, $P_c(t)$).
 - Important property: $F_{det}(id, P_c(t))$ approximates $F_{op}(id, S(t))$

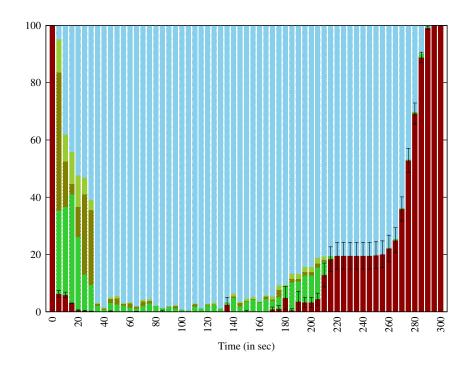
Can we observe the global structure? YES! Phase concept!!

- In case of Beernet, we can identify a qualitative property depending on neighbor behavior
- Phases of a node are analogous to solid, liquid and gaseous phases in physical system (e.g., water)
 - Solid: neighbors do not change (core ring).
 - Liquid: neighbors changing (branches).
 - Gaseous: no neighbors (isolated nodes).
- Three liquid sub-phases in terms of available functionalities and probability of facing an immediate phase transition.
 - liquid-1: if peer is on a branch with depth <= 2 and holds a stable finger table;
 - liquid-2: if peer is on a branch with depth > 2, but not tail of a branch. The finger table holds > 50% valid fingers;
 - liquid-3: if peer is on a branch with depth > 2, and it is tail of a branch. Most fingers are invalid or crashed.

Phase Transitions in SON: red, green and blue areas correspond to % of nodes on ring (*solid*), branches (*liquid*) and isolation (*gaseous*) respectively.



Increasing churn with time up to 100%, then decreasing churn with time:



What are Phase Transitions good for?

- ✓ Give useful information to the application.
- √ Can be used for efficient self-management.

Summary of Second Story

- The Phase of each node has a direct correlation with the overall functionalities (e.g., routing, availability of keys, transactions) of the system.
 - The current phase and phase transition at each node can be determined with high confidence, without any global synchronization.
- Reversible Phase Transitions in the system with varying stress can be observed as a by-product of making the system Reversible.
 - The system "boils" to the gaseous state (becomes disconnected) when churn increases and "condenses" from gaseous back to solid phase as churn intensity goes down.
 - Can provide useful information to the application layer using APIs.
 - Can be used for efficient self-management of the system.

Third Story Phase API and Applications

Phase API

An API exists on each node to expose its phase to the application layer

- Push and pull methods to communicate the current phase of a node
 - getPhase(? P_{cur}) Binds P_{cur} to the current phase of the peer.
 - setPhaseNotify(f) Sets a user-defined function, $f(?P_{new})$ to be executed when the phase changes. P_{new} is bound to the next phase of the peer and f is executed. Executions of f are serialized in the same thread over a stream of successive phases.

Phase-Aware Applications

- Predictable behavior for the users: an indicator that changes color to indicate the current phase of the underlying node.
 - Allow users to work productively offline and prevent any potential data-loss.
- Reversibility for the application:
 - Can increase replication factor of critical data, based on phase of underlying node;
 - Can improve throughput, by adapting philosophy of exponential back-off as TCP congestion algorithm.
 - Can manage its behavior for congestion-avoidance, thus help system to recover quickly.
- Empirical Demonstration of Phase-Aware Application design (future work)

Conclusion and Future Work

Conclusion

- In order to design provably correct decentralized networked systems, it is required to ensure their reversibility against stressful environments.
 - Build systems that are both predictable (hence, useful in practice)
 and reversible (hence, they survive)
- We define the concept of Reversibility to make precise what survival means in stressful environments
- We define the concept of Phase to allow applications to observe their stressful environment and act accordingly

Summary of Our Stories

- First Story: Repeated join and merger using Knowledge Base is required to achieve Reversibility against extremely high Churn
- Second Story: We observe Phases and Phase Transitions in the system as a by-product of making the system Reversible (give useful information to applications using APIs)
- Third Story: We introduce a Phase API to give useful information to applications and use it for phase-aware application design: predictable behavior and reversibility in the application-level semantics.

Future Work

Continuing the work directly:

- Deepen the analogy between phase in SONs and in physical systems;
- Design applications that take advantage of the Phase API to survive in extremely stressful environments;
- Gain more insights about the maintenance strategies.

Other topics:

- Investigate other application architectures;
- Investigate other stresses and stress interactions;
- Move to real environment, not simulated.

Thank You!!

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