

Documenting Evolving Software Systems through Reuse Contracts

Submitted to the OOPSLA'96 Workshop on "Object-Oriented Software Evolution and Re-Engineering"

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are not affected when the structures and implementations behind the interfaces are changed behind interfaces. As these other parts only rely on the information they get from these interfaces, they software engineering principle to cope with change and evolution. This principle is the foundation of interfering with the other parts. Details that are of no importance to other parts of the system are hidden about different system parts separately as well as making changes to certain parts of a system without amongst others, encapsulation, modularity, high cohesion and loose coupling. It enables reasoning Minimisation of dependencies between the parts of a software system is by far the most successful

made in software engineering, it can only take us so far. At a certain point in the evolution of a software system, changes occur that cannot be kept local to one system part and thus interfaces do have While the continuous elaboration on this principle accounts for much of the progress that has been

on where problems might occur upon change; the latter provides us information on what the problem is what other parts, but more importantly how they depend on each other. The former gives an indication between different system parts. Such a documentation must not only include which parts depend on development of software. This can only be dealt with by a careful documentation of dependencies (and thus on how it can be solved). The lack of this kind of documentation is a major impediment to Assessing the impact of such non-local changes remains one of the most compelling problems in the

abstraction intuitively correspond to the (partial) breaching of a contract. extension, refinement and concretisation. The inverse operators: coarsening, cancellation and and customised, reuse contracts are subject to parallel changes encoded by formal *reuse operators*: reuse contracts. Reuse contracts not only document how a system part can be reused, but also how and We propose to document the protocol between designers of different parts of a system by means of why the part is actually reused by other parts. Just as real world contracts can be extended, amended

Reuse contracts together with their operators facilitate managing the evolution of a software system by indicating how much work is needed to update the system, by forecasting when and which problems might occur, and by providing information on where and how to test and adjust the system.

Managing Parent Class Exchange

The use of abstract classes with inheritance as reuse mechanism is undoubtedly the best-known technique available today for structuring and adapting object-oriented software. Therefore, in [Steyaert&al96] we focused on the problem of evolution of class-hierarchies as a more tangible case to explain the ideas behind reuse contracts. In that context, reuse contracts and their operators describe the protocol between managers and users of (abstract) class libraries. Reuse contracts of abstract classes provide an explicit representation of the design decisions behind an abstract class, including information such as: which methods can be sent to the class, which methods are invoked by what other methods, which methods are abstract or concrete, relationships with other classes, ... Only information relevant to the design is included. For example, auxiliary or implementation-specific methods are not mentioned in a reuse contract.

Reuse contracts can be manipulated by means of reuse operators. Refinement refines the design of some methods, extension adds new methods, concretisation makes abstract methods concrete. These reuse operators not only allow documenting the changes (and the intentions of these changes) made to a class, but a careful investigation of their interactions also allows to predict and manage the effect of these changes.

Consider the example of a Collection hierarchy. A class Set defines a method add and a method addall to add a collection of elements simultaneously.

Class Set method add(Element) = 0 method addAll(aSet:Set) = begin for e in aSet do self.add(e) end

end

In order to decide which methods need to be overridden when creating a subclass CountableSet of Set that keeps a count of the number of elements in the set, we need information on which methods depend on what other methods. For example, if we know that addAll depends in its implementation on add, it is sufficient to override the method add to take counting into account. Reuse contracts for classes document these dependencies. In a reuse contract each method has a specialisation clause (in italics in the example) that documents how it depends on the other methods from this reuse contract (as in Lamping's specialisation interfaces [Lamping93]). The reuse contract is an interface description to

which the implementation must comply. It provides information that is typically not included in other

methodologies

reuse contract Set
abstract
add(Element)
concrete
addAll(Set) {add(Element)}
end

Reuse contracts in their current form only document the internal dependencies among a class's methods. Part of our future work is studying how reuse contracts can be extended to include interclass dependencies as well. Other experiments are being conducted to include information on state and state transformations in reuse contracts and operators.

Reuse contracts can be used to assess the impact of changes or updates to system parts. Suppose we want to make an optimised version OptimisedSet of Set. In this version addAll stores the added elements directly rather than invoking the add method to do this. This leads to inconsistent behaviour in CountableSet when it decides to upgrade OptimisedSet; not all additions will be counted. This is because the assumption made by CountableSet that addAll invokes add is broken in OptimisedSet. Using the terminology of [Kiczales&Lamping92] we say that addAll and add have become *inconsistent methods*. Although in this simple example the conflict can easily be derived from the code, in practice it should be possible to detect such conflicts without code inspection. The major obstacle for locating problems such as inconsistent methods is that the different conceptual ways to reuse an (abstract) class are all performed by the same operator, i.e. inheritance. More information about the intentions of inheritor is needed. This kind of information is precisely provided by the reuse operators on reuse contracts. In the example the reuse contract of CountableSet and OptimisedSet document how they were derived from Set.

reuse contract CountableSet concretises Set concrete add(Element) end

use contract Optim

reuse contract OptimisedSet coarsens Set concrete addAll(Element) {-add(Element})

The fact that add and addAll have become inconsistent can be derived directly from the reuse contracts: CountableSet concretises a method that has been removed from the specialisation clause while changing from the old parent class to the new parent class (in italics above).

For a more complete set of possible conflicts on parent class exchange we refer to [Steyaert&al96]. Furthermore, it gives a complete set of reuse operators together with a set of rules that allow automatic detection of conflicts based on the interaction of reuse operators.

Environment and Tool Support for Reuse Contracts	
An environment for managing software evolution based on the concept of reuse contracts should	
include tool support for assessing the impact of making changes to a system by signalling possible problems that (might) occur. A prototype version of such a tool has been implemented in PROLOG.	
The environment can also assist in the synchronisation of reuse contracts and their corresponding	
lesign, the implementation must be forced to comply to the reuse contract. In those parts that are still	
subject to major redesign, it should be possible to make changes to both implementation and reuse	
contracts independently. The environment could discretely issue warnings, but should not become a nindrance.	
Finally, for software systems that have not been documented by means of reuse contracts, tools can be	
constructed that semi-automatically extract this documentation from the code, based on the calling	
locumentation, as reuse contracts should include only information relevant to the design. Once the	
lifferent reuse contracts have been extracted, the tool can easily compute how the reuse contracts	
corresponding to the different parts of the system are related to one another by means of reuse	
Conclusion	
Current methodological and tool support for managing the evolution of large, long-lived software	
ystems, focuses mainly on minimising dependencies between system parts. However, the question	
what happens when these dependencies are changed at some point during the evolution process is argely neglected. Documenting these dependencies by means of reuse contracts and reuse operators	
allows us to signal such changes and to assess their impact. Many tools to support the use of reuse	
contracts for managing software evolution can be conceived. When adopted, reuse contracts may	
significantly enhance the way in which software is being built and managed.	
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